scrypt: A new key derivation function
Doing our best to thwart TLAs armed with ASICs

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Making bcrypt obsolete

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Are you sure your SSH keys are safe?

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What are key derivation functions?

- You have a password.
- You want a generate a *derived key* from that password.
  - Verifying passwords for user authentication.
  - Encrypting or signing files.
- In most situations where passwords are used, they are passed to a key derivation function first.
  - In most situations where key derivation functions aren’t used, they should be!
- **Examples of key derivation functions:**
  - DES CRYPT [R. Morris, 1979]
  - MD5 CRYPT [P. H. Kamp, 1994]
  - bcrypt [N. Provos and D. Mazières, 1999]
  - PBKDF2 [B. Kaliski, 2000]
  - MD5 (not really a key derivation function!)
Attack model: Assume that the attacker can mount an offline attack.

- Attacker has access to /etc/master.passwd and wants to find the users’ passwords.
- Attacker has an encrypted file and wants to decrypt it.

For all reasonable key derivation functions, the only feasible attack is to repeatedly try passwords until you find the right one.

- This is called a “brute force” attack.

If it takes twice as long to check if a password is correct, it will take twice as long to find the right password.

- ...as long as the attacker is using the same software as you.
Hardware-based brute force attacks

- Some organizations have the resources to design and fabricate custom password-cracking integrated circuits.
  - The US National Security Agency
  - The UK Government Communications Headquarters?
  - The Communications Security Establishment of Canada?
  - The Chinese government?
  - Organized crime?
  - The Electronic Frontier Foundation?

- Using ASICs, it is possible to pack many copies of a cryptographic circuit onto a single piece of silicon.

- Moore’s law: Every 18–24 months, a new generation of semiconductor manufacturing processes makes CPUs faster.
  - …password-cracking ASICs get faster AND can fit more copies of a password-cracking circuit.
The cost of a hardware brute-force attack is measured in dollar-seconds.

Password cracking is embarrassingly parallel, so if you use twice as much hardware you can crack the key in half the time.

Cost of ASICs \(\propto\) size of ASICs.

A strong key derivation function is one which can only be computed by using a large circuit for a long time.

J. Kelsey, B. Schneier, C. Hall and D. Wagner, 1998: Use “32-bit arithmetic and moderately large amounts of RAM”.

An example of a “moderately large amount of RAM”: 1 kB.

If we use a *ridiculously* large amount of RAM, hardware attacks will be even more expensive.
Memory-hard algorithms

Definition

A *memory-hard* algorithm on a Random Access Machine is an algorithm which uses $S(n)$ space and $T(n)$ operations, where $S(n) \in \Omega \left( T(n)^{1-\epsilon} \right)$.

- Conceptually, a memory-hard algorithm is one which comes close to using the largest amount of storage possible for an algorithm with the same running time. 
  
  ... and consequently the largest circuit area possible.

- The HEKS key derivation algorithm [A.G. Reinhold, 1999] is memory-hard, but it isn’t very secure, since it can be effectively parallelized.

- Secure key derivation functions require a large die area *and* a lot of time to compute.

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A sequential memory-hard function is a function which (a) can be computed by a memory-hard algorithm on a Random Access Machine in $T(n)$ operations; and (b) cannot be computed on a Parallel Random Access Machine with $S^*(n)$ processors and $S^*(n)$ space in expected time $T^*(n)$ where $S^*(n)T^*(n) = O(T(n)^{2-x})$ for any $x > 0$.

- Not only do memory-hard functions require lots of storage, but they also cannot be parallelized efficiently.
- If we can find a key derivation function which is sequential memory-hard, it should be very secure against hardware attack.
Algorithm (ROMix)

Given a hash function $H$, an input $B$, and an integer parameter $N$, compute

$$V_i = H^i(B) \quad 0 \leq i < N$$

and $X = H^N(B)$, then iterate

$$j \leftarrow \text{Integerify}(X) \mod N$$
$$X \leftarrow H(X \oplus V_j)$$

$N$ times; and output $X$.

- The function $\text{Integerify}$ can be any bijection from $\{0, 1\}^k$ to $\{0 \ldots 2^k - 1\}$.
- ROMix fills $V$ with pseudorandom values, then accesses them in a pseudorandom order.
Theorem

Under the random oracle model, the class of functions ROMix are sequential memory-hard.

- The random oracle model is a very standard assumption in proofs relating to hash functions.
- The proof is roughly 2 pages long.
  - I could probably spend my entire talk explaining the proof.
  - ... but I won’t.
  - If you’re interested, go read the paper I wrote about this.
- It’s much easier to prove that an algorithm runs in specified time and space than to prove a minimum bound on the time and space used by any algorithm which computes a function.
Use PBKDF2 to convert a password into a bitstream.

Feed this bitstream to ROMix.

Feed the output of ROMix back to PBKDF2 to generate the derived key.

Cryptographic primitives used:
  - HMAC-SHA256
  - Salsa20/8 core

The Salsa20/8 core outputs lots of bits very fast, which means that scrypt can use lots of memory.
  - Approximately 1 byte of RAM per 10 clock cycles on a Core 2 processor.
  - We can quickly require a large semiconductor area.
It's hard to get accurate information about how much it costs to build password-cracking machines.

Oddly enough, the NSA doesn’t publish this data.

The best we can do for most KDFs is to count cryptographic operations and assume that they are responsible for most of the time and die area.

This is probably a fairly accurate approximation, since key derivation functions only have a very small amount of non-cryptographic computations.

For scrypt we also need to look at the die area required for storage.
Approximate circuit complexity and performance for cryptographic primitives, based on a 130 nm semiconductor process:

- A DES circuit with \( \approx 4000 \) gates of logic can encrypt data at 2000 Mbps.
- An MD5 circuit with \( \approx 12000 \) gates of logic can hash data at 2500 Mbps.
- A SHA256 circuit with \( \approx 20000 \) gates of logic can hash data at 2500 Mbps.
- A Blowfish circuit with \( \approx 22000 \) gates of logic and 4 kiB of SRAM can encrypt data at 1000 Mbps.
- A Salsa20/8 circuit with \( \approx 24000 \) gates of logic can output a keystream at 2000 Mbps.

I’m using 130 nm as a reference point because this is what I could get the most data for.
Estimating hardware brute force attack costs

- Very approximate estimates of VLSI area and cost:
  - Each gate of random logic requires $\approx 5 \mu m^2$ of VLSI area.
  - Each bit of SRAM requires $\approx 2.5 \mu m^2$ of VLSI area.
  - Each bit of DRAM requires $\approx 0.1 \mu m^2$ of VLSI area.
  - VLSI circuits cost $\approx 0.1$/mm$^2$.

- These values are based on a 130 nm process circa 2002.
- These values have a very wide error margin.
  - Non-cryptographic parts of ASICs (e.g., I/O), chip packaging, boards, power supplies, and operating costs could increase password-cracking costs by a factor of 10.
  - Improvements in semiconductor technology since 2002 could reduce password-cracking costs by a factor of 10.
  - Improved cryptographic circuits could reduce costs by a factor of 10.
Key derivation functions

- Non-parameterized KDFs:
  - DES CRYPT
  - MD5 CRYPT
  - MD5

- KDFs tuned for interactive logins \((t \leq 100 \text{ ms})\):
  - PBKDF2-HMAC-SHA256, \(c = 86000\)
  - bcrypt, \(cost = 11\)
  - scrypt, \(N = 2^{14}, r = 8, p = 1\)

- KDFs tuned for file encryption \((t \leq 5 \text{ s})\):
  - PBKDF2-HMAC-SHA256, \(c = 4300000\)
  - bcrypt, \(cost = 16\)
  - scrypt, \(N = 2^{20}, r = 8, p = 1\)

- Running time based on a 2.5 GHz Core 2 (aka. my laptop).
Passwords

- 6 lower-case letters; e.g., “sfgroy”.
- 8 lower-case letters; e.g., “ksuvnwyf”.
- 8 characters selected from the 95 printable 7-bit ASCII characters; e.g., “6,uh3y[a”.
- 10 characters selected from the 95 printable 7-bit ASCII characters; e.g., “H.*W8Jz&r3”.
- A 40-character string of text; e.g., “This is a 40-character string of English”.
  - Entropy estimated according to formula from NIST: 1st character has 4 bits of entropy; 2nd–8th characters have 2 bits of entropy each; 9th–20th characters have 1.5 bits of entropy each; 21st and later characters have 1 bit of entropy each.
Estimated brute force attack costs

Estimated cost of hardware to crack a password in 1 year.

<table>
<thead>
<tr>
<th>KDF</th>
<th>6 letters</th>
<th>8 letters</th>
<th>8 chars</th>
<th>10 chars</th>
<th>40-char text</th>
</tr>
</thead>
<tbody>
<tr>
<td>DES CRYPT</td>
<td>&lt; $1</td>
<td>&lt; $1</td>
<td>&lt; $1</td>
<td>&lt; $1</td>
<td>&lt; $1</td>
</tr>
<tr>
<td>MD5</td>
<td>&lt; $1</td>
<td>&lt; $1</td>
<td>&lt; $1</td>
<td>$1.1k</td>
<td>$1</td>
</tr>
<tr>
<td>MD5 CRYPT</td>
<td>&lt; $1</td>
<td>&lt; $1</td>
<td>$130</td>
<td>$1.1M</td>
<td>$1.4k</td>
</tr>
<tr>
<td>PBKDF2 (100 ms)</td>
<td>&lt; $1</td>
<td>&lt; $1</td>
<td>$18k</td>
<td>$160M</td>
<td>$200k</td>
</tr>
<tr>
<td>bcrypt (95 ms)</td>
<td>&lt; $1</td>
<td>$4</td>
<td>$130k</td>
<td>$1.2B</td>
<td>$1.5M</td>
</tr>
<tr>
<td>scrypt (64 ms)</td>
<td>&lt; $1</td>
<td>$150</td>
<td>$4.8M</td>
<td>$43B</td>
<td>$52M</td>
</tr>
<tr>
<td>PBKDF2 (5.0 s)</td>
<td>&lt; $1</td>
<td>$29</td>
<td>$920k</td>
<td>$8.3B</td>
<td>$10M</td>
</tr>
<tr>
<td>bcrypt (3.0 s)</td>
<td>&lt; $1</td>
<td>$130</td>
<td>$4.3M</td>
<td>$39B</td>
<td>$47M</td>
</tr>
<tr>
<td>scrypt (3.8 s)</td>
<td>$900</td>
<td>$610k</td>
<td>$19B</td>
<td>$175T</td>
<td>$210B</td>
</tr>
</tbody>
</table>

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scrypt: A new key derivation function
When used for interactive logins, scrypt is . . .
- \( \approx 2^5 \) times more expensive to attack than bcrypt,
- \( \approx 2^8 \) times more expensive to attack than PBKDF2,
- and \( \approx 2^{15} \) times more expensive to attack than MD5 CRYPT.

When used for file encryption, scrypt is . . .
- \( \approx 2^{12} \) times more expensive to attack than bcrypt,
- \( \approx 2^{15} \) times more expensive to attack than PBKDF2,
- and \( \approx 2^{37} \) times more expensive to attack than MD5.

openSSL enc uses MD5 as a key derivation function.

OpenSSH uses MD5 as a key derivation function for passphrases on key files.
- Are you sure that your SSH keys are safe?
More details at http://www.tarsnap.com/scrypt/.

- Source code for scrypt.
- A simple file encryption/decryption utility.
- A 16-page paper.

Questions?