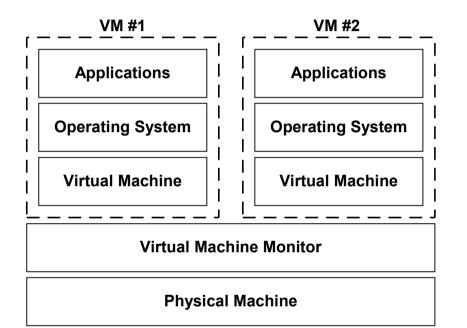
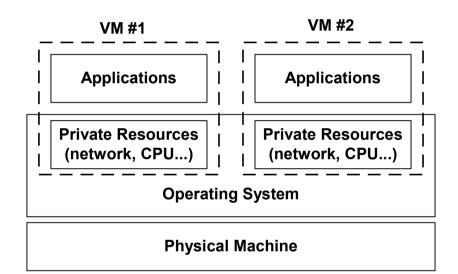
# Network stack virtualization for FreeBSD 7.0

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- Network stack virtualization what, why, and how?
- Who needs this?
- Implementation: FreeBSD 4.x vs. 7.0
- Generalizing OS-level resource virtualization?

# Server virtualization: two sides of the spectrum





Strong isolation model Independent OS instances VM migration possible Efficient resource utilization

No extra I/O overhead

Scaling

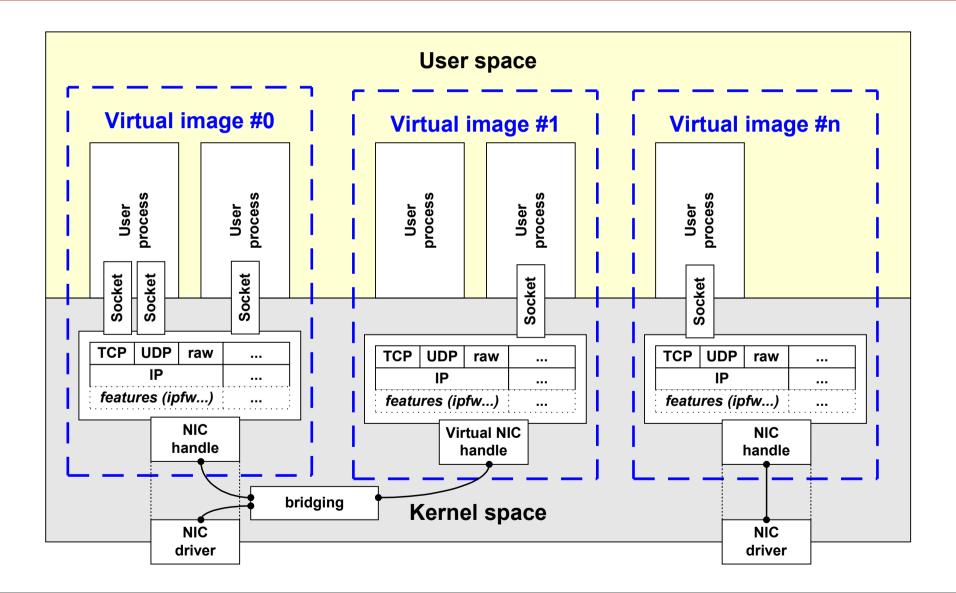
- Traditional OS architecture
  - Support for only a single instance of network stack or protocol family within the kernel
  - Jails: first successful pseudo-virtualization framework
- Network stack virtualization (or *cloning*)
  - Multiple independent network stack <u>state</u> instances within a single kernel
  - Existing networking code paths and algorithms remain the same, but must be taught on how to operate on virtualized symbols / state

# Applications: who needs this?

#### - Virtual hosting

- Think of extending FreeBSD jail with its own independent network stack instance: multiple interfaces and IP addresses, private routing table, IPFW / PF, dummynet, BPF, raw sockets etc. etc.
- Anecdotal evidence: FreeBSD 4.11 based version in production use by some US ISPs
- VPN provisioning and monitoring
  - Support for overlapping IP addressing schemes
- Network simulation / emulation
  - Each network stack instance == an independent virtual node or router -><u>http://www.imunes.net/</u>

# The basic idea: replicate global networking state



# Implementation concepts: long time ago...

- Patches against FreeBSD 4.7 .. 4.11 kernels
  - Obsolete platform today
- struct vnet
  - One huge structure / container; each network stack instance operates on its private copy
  - Contains ifnet lists, IPv4 / IPv6 / firewall state etc.

# Sockets

- Each socket is assigned to a network stack instance during creation time
- Cannot move / change until socket closed

Implementation concepts: how it was done (cont'd)

- Network interfaces
  - Each interface can belong to only one network stack instance at a time
  - Demultiplexing of incoming traffic based on on the new if\_vnet tag in struct ifnet
  - Network communication between stack instances only through explicit links: netgraph
- User processes
  - Bound to only one stack at a time, can reassociate
  - Jail-style separation (reused existing jail code)

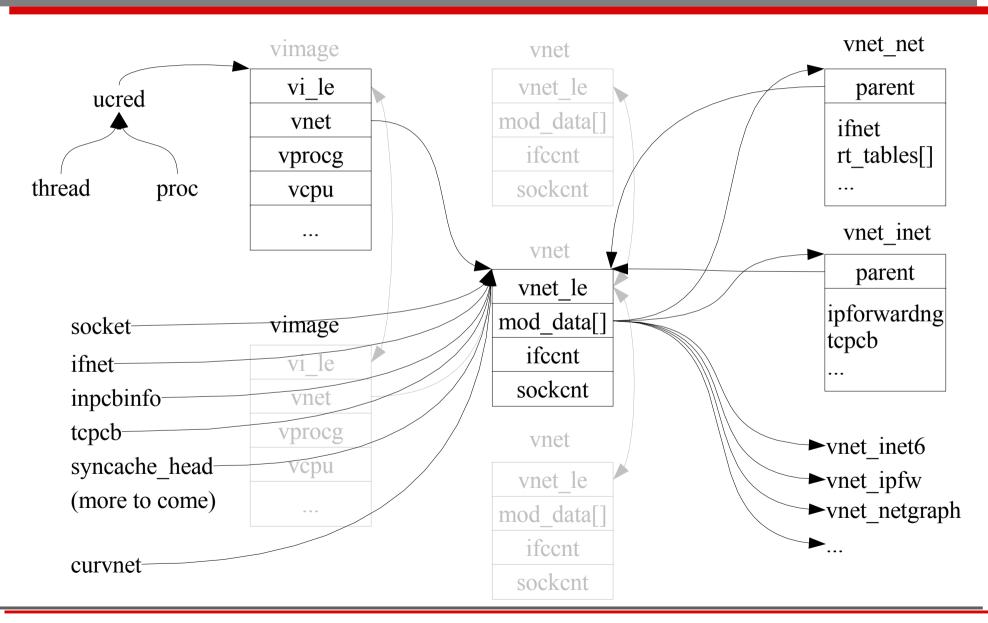
# Implementation concepts: API / ABI compatibility

- Userland to kernel: both API and ABI 100% preserved
  - Support for accessing the virtualized symbols added to the kldsym interface (needed for netstat, systat, top and similar utilities)
  - Similar extensions added to the <code>sysctl</code> interface
- Within the kernel: API is *NOT* preserved
  - Many networking functions extended with an additional argument: struct vnet \*
  - Generally, no changes at device driver layer

# (Re)implementation: 7.0

- Goals:
- Conditional compilation
- Better support for kernel loadable modules
- Scope of changes is huge: reduce code churn
- SMP must work
- Otherwise, no chances for including the changes into main FreeBSD tree

# Replicate global networking state: how?



#### vnet modules: registration / deregistration

```
static struct vnet symmap vnet net symmap[] = {
       VNET SYMMAP(net, ifnet),
       VNET SYMMAP(net, rt tables),
       VNET SYMMAP END
};
static struct vnet modinfo vnet net modinfo = {
        .id
                       = VNET MOD NET,
        .flaqs
                       = VNET MFLAG ORDER 1ST,
                       = "net",
        .name
        .symmap = vnet_net_symmap,
        .i attach = vnet net iattach,
        .i detach
                       = vnet net idetach
};
if init(void *dummy unused)
#ifdef VIMAGE
       vnet_mod_register(&vnet_net_modinfo);
#else
       vnet net iattach();
#endif
```

## **Conditional compilation:** option VIMAGE

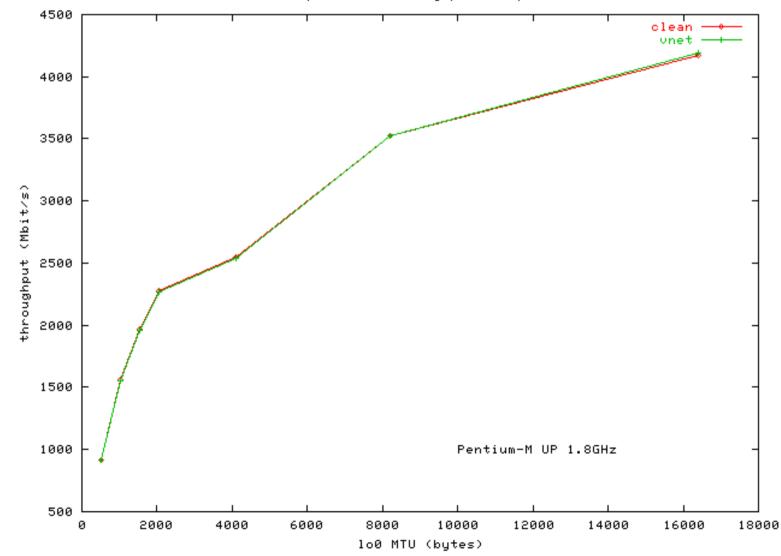
- Dereference virtualized symbols: how?
  - Use macros for this. Example:
    - if\_addrhead **becomes v\_**if\_addrhead
  - Standard kernel:
    - V\_if\_addrhead expands back to if\_addrhead
  - Virtualized kernel:
    - V\_if\_addrhead expands to vnet\_net->\_if\_addrhead
  - Sysctl and kldsym interfaces extended to support access to virtualized symbols

- Implicitly pass the vnet context to operate on:
  - Thread-local curvnet variable

```
void if_attach(struct ifnet *ifp)
{
    INIT_VNET_NET(curvnet);
    ...
```

```
INIT_VNET_NET(x) (x is a struct vnet *) expands to
struct vnet_net *vnet_net = x->mod_data[VNET_MOD_NET];
```

# Performance: loopback TCP throughput



Loopback TCP throughput (netperf)

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# Generalizing OS-level virtualization

- Management concepts / API
  - Top-level resource container struct vimage
  - Contains freely combinable subsystem-specific state
    - vnet, vcpu, vprocg, vfs...
  - Single process with sockets in multiple stacks
    - Extend socket interface -> multi-table routing daemons
  - Hierarchy of vimages follow UNIX process model?
  - Permissions, restrictions, inheritance...
  - How to best integrate those new concepts / features with the rest of the system?

# Project status

- Supported by NLNet and FreeBSD foundation
  - Started in August 2006, should have already finished...
- In sync with -CURRENT: p4 projects/vimage
  - Snap-in replacement kernel no userspace changes!
  - <u>http://imunes.tel.fer.hr/virtnet/</u> : CVSup
- Reasonably stable already
  - Lots to be done: locking, management API & housekeeping
- Most important networking subsystems virtualized:
  - IPv4, IPv6, NFS, IPFW / PF firewalls, BPF, raw / routing sockets...
- Outside the tree until 7.0-RELEASE, merging in 8.0?

# To conclude...

- Do we need all this?
  - the community has to provide that answer.
- If yes, what's next to virtualize?
  - CPU time (scheduler)
  - Filesystems (ZFS?) / disk I/O bandwidth
  - Memory
  - ...
- We need a generalized OS-level virtualization model

#### http://imunes.tel.fer.hr/virtnet/